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# **Modelling Go Positions with Planar CRFs**

## **Statistical Machine Learning Program**

### Abstract

We apply Conditional Random Fields (CRFs) to territory prediction in the game of Go. We propose a two-stage graph reduction of the Go position, with the first stage used for CRF parameter estimation and the second for inference. The interaction potentials in our model are calculated from generic shape features; the Figure 1. associated parameters are shared between semantically equivalent feature types. Our experiments indicate that this architecture is very efficient at propagating relevant information across the graph.

### Go

- Two players (black and white) alternate in placing stones on the intersections of a  $9 \times 9$  (or  $19 \times 19$ ) grid.
- Neighbouring stones of the same colour form a contiguous *block*.
- A block can be *captured* if all its empty neighbours (*liberties*) are occupied by opponent stones.
- Players aim to maximize the area (*territory*) of blocks that cannot be captured.
- *Territory prediction*: given a board position, determine which player controls each intersection (see Figure 1 for an example).

### **Graph Abstraction**

### Common fate graph

- Common fate property: blocks always live or a die as a unit.
- Common fate graph  $G_f$  (Graepel et al., 2001) merges all black stones in a block into a single node  $\bullet$  and white stones into  $\bigcirc$  (Figure 2).
- Block's size and shape can be represented in the corresponding node's features.

### **Block graph**

- Empty regions are usually divided between players and so cannot be fully collapsed. However if we do not collapse them at all, we end up with large, unwieldy graphs.
- Our *block graph*  $G_b$  is a compromise between these two extremes (Figure 3).
- We collapse empty intersections of  $G_f$  into black surround  $\blacksquare$ , white surround  $\Box$  and *neutral*  $\diamond$  depending on the Manhattan distance to the nearest black and white stones.
- $G_b$  is more concise than  $G_f$ , but preserves all the information required for predicting territory.
- Since  $G_h$  is planar, we employ exact polynomial-time parameter estimation based on the work of Globerson and Jaakkola (2007).

### Group graph

- *Group*: set of blocks of the same colour that share at least one surround.
- We construct the group graph  $G_g$  by collapsing groups of  $G_b$  (Figure 3).
- Blocks belonging to the same group are likely to have the same fate, therefore we use  $G_g$  for inference.

Figure 2. Corresponding common fate graph  $G_f$ .

Figure 3. Corresponding block graph  $G_b$ .  $\bigcirc$ represent *stones*, □ represent surrounds, and ♦ are *neutral*. Dashed lines indicate nodes of the group graph  $G_g$ .

A typical  $9 \times 9$  board position represented as a grid graph G. Black's territory is shaded in darkgray, while White's territory is shaded in lightgray. The stones inside the opponent's territory are dead; the rest are alive.



Example

11

14

12 + 13

- 16 - 17

19

- 4 measures of error



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### **Features and Parameters**

### Nodes

• Given a node  $v \in G_b$ , for each point  $i \in v$  compute the number of adjacent points  $A_i$  that are also in v (see Figure 4 top).

• Node's feature is a vector *F*, where  $F_k = |\{i : A_i = k\}|$ .

• F provides a powerful summary of the region's shape. This is important for classifying the region's eye-forming potential and thus its ability to live/die.

### Edges

• Given two adjacent nodes  $v_1, v_2 \in G_b$ , for each point  $i \in v_1$  compute the number of adjacent points  $A_i^1$  that are in  $v_2$  and vice versa for  $A_i^2$ .

• Edge's features are two vectors  $F^1$  and  $F^2$  that use  $A^1$  and  $A^2$  respectively. • Our edge features provide additional information that is not conveyed by the features of the nodes they connect (see Figure 4 bottom).

### Parameter sharing

• We encode colour symmetry by appropriate parameter sharing.

• 3 node parameters: *stone* ( $\theta_{\bigcirc}$ ), *surround* ( $\theta_{\square}$ ), and *neutral* ( $\theta_{\diamond}$ ).

• 8 edge parameters: one for each possible type of node pairing.

• For each edge we also include features of its neighbouring nodes and edges. These features are associated with the parameter vector  $\theta^n$ .

### Experiments

• We use  $9 \times 9$  endgame positions of van der Werf et al. (2005). • Each intersection is labeled as either BLACK OF WHITE.

- Vertex: percentage of misclassified *non-neutral* nodes in G.

**– Block**: percentage of misclassified *stone* nodes in  $G_h$ .

- Winner: percentage of games whose winner is predicted incorrectly. - Game: percentage of games that have at least one vertex error.

• Naive error assumes all stones are alive, *i.e.*,  $\bullet$  is BLACK and  $\bigcirc$  is white.

### **Conclusion and Future Work**

• Our system is clearly outperforming previous generic models, however it is worse than approaches that incorporate Go-specific features.

• Our high-order graph representation enables the system to capture the subtleties of life and death in Go, despite our simple generic features.

• Future work will focus on better ways of classifying empty regions and incorporation of more domain-specific knowledge.

Figure 4. Top: Node features, resulting in feature vector [2, 4, 2, 1]. Bottom: Edge features, resulting in feature vectors [6, 3, 0] and [3, 3, 1] for edges  $\bullet \to \bigcirc$ and  $\bigcirc \rightarrow \bullet$ , respectively.

Figure 5. Paramet tures used to compu tial of one particular in a small block gra





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		Current Edge		Neighbour Edges	
ters and fea- ute the poten- ar edge (blue) aph (below).		Param.	Feat.	Param.	Feat.
	Nodes	$ heta_{\bigcirc}$		$oldsymbol{ heta}^n_\diamond$	\$
		$ heta_{\Box}$	•	$ heta_{\bigcirc}^n$	0
				$ heta_{\Box}^{n}$	
	Edges	$ heta_{\bigcirc\Box}$	$\bullet \to \blacksquare$	$ heta_{\diamond \bigcirc}^n$	$\diamond \rightarrow \bullet$
		$ heta_{\Box \bigcirc}$	$\blacksquare \to \blacklozenge$	$ heta_{\bigcirc\diamond}^n$	$\bullet \to \diamond$
				$ heta_{\bigcirc\bigcirc}^n$	$\bigcirc \rightarrow ullet$
					$\bullet \to \bigcirc$
				$ heta_{\Box\Box}^n$	$\Box \rightarrow \blacksquare$
					$\blacksquare \to \square$

	Error (%)					
Algorithm	Vertex	Block	Winner	Game		
Naive	6.79	17.57	30.79	75.70		
rn et al. (2004)	4.77	7.36	13.80	38.30		
Block graph	2.36	3.56	4.53	13.02		
h + neighbour features	1.87	2.76	3.42	9.60		
+ other enhancements	1.54	2.20	2.09	7.90		
* GnuGo	-	-	-	1.32		
er Werf et al. (2005)	0.19	$\leq 1.00$	0.50	1.10		

Figure 6. Prediction error (percentage) for various algorithms. \*: employs Go- specific features and was used to label training data.

### References

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