

SPANNING TREES IN RANDOM REGULAR GRAPHS

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Let $n_1 < n_2 < \dots$ be possible orders of connected regular graphs of fixed degree $k \geq 3$. For each i , choose a graph X_i at random from the set of all connected regular graphs of order n_i and degree k . Let $\kappa(X_i)$ be the number of spanning trees of X_i . Then, with probability one,

$$\kappa(X_i)^{1/n_i} \rightarrow \frac{(k-1)^{k-1}}{(k^2-2k)^{k/2-1}} \quad \text{as } i \rightarrow \infty.$$

§1 INTRODUCTION

Let $X = X_1, X_2, \dots$ be a sequence of connected regular graphs of degree $k \geq 3$ and orders $n_1 < n_2 < \dots$. Define κ_i to be the number of spanning trees of X_i . For $i \geq 1$ and $m \geq 3$, let $C(m, X_i)$ be the number of cycles of length at most m in X_i . We will say that X satisfies Condition (B) if there are constants m_0 (arbitrary), $m_1 > 1/\log(k^2/4(k-1))$ and $\epsilon > 0$ such that $C(2\lfloor m_0 + m_1 \log \log n_i \rfloor, X_i) = O(n_i(\log n_i)^{-1-\epsilon})$ as $i \rightarrow \infty$. The ordinal B is used for consistency with [2]. Our first theorem was proved in [2].

1.1 Theorem *If X satisfies Condition (B), then*

$$\kappa(X_i)^{1/n_i} \rightarrow \frac{(k-1)^{k-1}}{(k^2-2k)^{k/2-1}} \quad \text{as } i \rightarrow \infty. \quad \square$$

In this paper we will prove that, if X is constructed by choosing each X_i at random from the set of all connected labelled regular graphs of order n_i and degree k , then X satisfies Condition (B) with probability one.

§2 SUBGRAPHS OF RANDOM REGULAR GRAPHS

In this section we will be concerned exclusively with the set $R(n, k)$ of all connected labelled regular graphs of degree k and order n . Two such graphs, X_1 and X_2 , are *closely related* if $|E(X_1) \setminus E(X_2)| = 2$. Since we are only considering simple graphs, it is clear that each of $E(X_1) \setminus E(X_2)$ and $E(X_2) \setminus E(X_1)$ consists of two independent edges and that these four edges together form a square. Another way of defining this relationship is to say that $X_1, X_2 \in R(n, k)$ are closely related if there are distinct vertices v_1, v_2, v_3, v_4 of X_1 such that $v_1v_2, v_3v_4 \in E(X_1)$, $v_1v_3, v_2v_4 \notin E(X_1)$ and $E(X_2) = E(X_1) \setminus \{v_1v_2, v_3v_4\} \cup \{v_1v_3, v_2v_4\}$.

Let F be any non-empty subset of $E(K_n)$, and define $R(n, k, F)$ to be the subset of $R(n, k)$ consisting of those graphs X for which $F \subseteq E(X)$. Let $v_1v_2 \in F$.

2.1 Lemma $|R(n, k, F \setminus \{v_1v_2\})| \geq \frac{nk - 2|F| - 2(k^2 - 1)}{2k^2} |R(n, k, F)|$.

Proof: Let M be the number of pairs (X_1, X_2) of closely related graphs, where $X_1 \in R(n, k, F)$ and $X_2 \in R(n, k, F \setminus \{v_1v_2\})$. We will prove the lemma by estimating M in two different ways.

Firstly, consider an arbitrary $X_1 \in R(n, k, F)$. We can construct any closely related graph in $R(n, k, F \setminus \{v_1v_2\})$ by finding adjacent vertices v_3 and v_4 of X_1 and then replacing the edges v_1v_2 and v_3v_4 by v_1v_3 and v_2v_4 . For this to be valid, it is necessary that $\{v_3, v_4\} \cap \{v_1, v_2\} = \emptyset$, $v_3v_4 \notin F$, and $v_1v_3, v_2v_4 \notin E(X_1)$. This leaves us with at least $nk - 2|F| - 2(k^2 - 1)$ choices for (v_3, v_4) , except that we are also requiring X_2 to be connected. Furthermore, it is easy to see that, since X_1 is connected, at least one of X_2 and X_2' is connected, where $E(X_2) = E(X_1) \setminus \{v_1v_2, v_3v_4\} \cup \{v_1v_3, v_2v_4\}$ and $E(X_2') = E(X_1) \setminus \{v_1v_2, v_3v_4\} \cup \{v_1v_4, v_2v_3\}$. Therefore,

$$M \geq (nk/2 - |F| - k^2 + 1)|R(n, k, F)|. \quad (1)$$

Now consider an arbitrary $X_2 \in R(n, k, F \setminus \{v_1v_2\})$. Since X_2 has degree k , the number of closely related graphs in $R(n, k, F)$ is at most k^2 . Therefore

$$M \leq k^2|R(n, k, F \setminus \{v_1v_2\})|. \quad (2)$$

The lemma now follows easily from (1) and (2). \square

2.2 Corollary If $R(n, k) \neq \emptyset$ and $2|F| < nk - 2(k^2 - 1)$, then

$$\frac{|R(n, k, F)|}{|R(n, k)|} \leq \left(\frac{2k^2}{nk - 2(k^2 - 1) - 2|F|} \right)^{|F|}$$

§3 CYCLES

Our primary aim in this section is to prove that X satisfies Condition (B) with probability one. In doing so we will prove rather more.

Since a direct application of Corollary 2.2 to the estimation of $C(m, X_i)$ does not appear to be sufficient, we will instead bound the expectation of $C(m, X_i)^2$. The following lemma will prove very useful.

3.1 Lemma Let $1 \leq l < r \leq s \leq n$. The number of ways of choosing an r -cycle C_1 and s -cycle C_2 from K_n such that $|E(C_1) \cap E(C_2)| = l$ is

$$w(n, s, r, l) \leq \frac{n^{r+s-l-1}}{2} \left(1 + \sqrt{\frac{2(l-1)}{n}} \right)^{r+s-2l-2}$$

Proof: Since $l < r \leq s$, $E(C_1) \cap E(C_2)$ consists of a collection of disjoint paths. Suppose that the number of paths is t .

The number of ways of choosing t paths, containing l edges altogether, is

$$\binom{l-1}{t-1} \frac{n!}{2^t t! (n-l-t)!} < \frac{n^{l+t} (l-1)^{t-1}}{2^t t! (t-1)!}.$$

Given $E(C_1) \cap E(C_2)$, we can construct C_1 by arranging the t paths in a cycle and then inserting $r-l-t$ new vertices. The number of ways this can be done is

$$2^{t-1} (r-l-1)! \binom{n-l-t}{r-l-t} < 2^{t-1} n^{r-l-t} \frac{(r-l-1)!}{(r-l-t)!}.$$

Constructing C_2 similarly, we find that the pair (C_1, C_2) can be chosen in less than

$$\frac{n^{r+s-l-t} 2^{t-2}}{t} \binom{r-l-1}{t-1} \binom{s-l-1}{t-1}$$

ways.

The sum of this expression over t is bounded by the function given in the statement of the lemma. \square

We will use $\mathcal{E}(Z)$ to denote the expectation of a random variable Z .

3.2 Theorem Let $m = m(n) = O(n^{\frac{1}{2}-\epsilon})$, for some $\epsilon > 0$. Then

$$\mathcal{E}(C(m(n_i), X_i)^2) = O\left(\frac{(2k)^{2m}}{m^2}\right).$$

Proof: First note that, by Corollary 2.2, the probability that X_i contains any specified set of $t \leq m(n_i)$ edges is bounded above by $A(2k/n_i)^t$, for some constant A .

Now define F_1, F_2, \dots, F_N to be the set of all cycle (read sets of edges forming cycles) of length at most m in K_n . For $1 \leq j \leq N$, define

$$Z_j = \begin{cases} 0, & \text{if } F_j \not\subseteq E(X_i), \\ 1, & \text{if } F_j \subseteq E(X_i). \end{cases}$$

Then clearly

$$\mathcal{E}(C(m, X_i)^2) = \sum_{1 \leq u, v \leq N} \mathcal{E}(Z_u Z_v). \quad (3)$$

We will break the sum above into three parts, according to $|F_u|$, $|F_v|$ and $|F_u \cap F_v|$. Put $n = n_i$.

(a) First consider the contribution to (3) with $u = v$. This is clearly equal to

$$\begin{aligned} \sum_{1 \leq u \leq N} \mathcal{E}(Z_u) &\leq A \sum_{s=3}^m \frac{n!}{2s(n-s)!} \left(\frac{2k}{n}\right)^s \\ &< A \sum_{s=3}^m \frac{(2k)^s}{2s} \\ &= O\left(\frac{(2k)^m}{m}\right). \end{aligned}$$

(b) Next, consider the terms of (3) for which $|F_u \cap F_v| = 0$. The contribution here is

$$\begin{aligned} A \sum_{r=3}^m \sum_{s=3}^m \frac{n!}{2r(n-r)!} \frac{n!}{2s(n-s)!} \left(\frac{2k}{n}\right)^{r+s} &< A \sum_{r=3}^m \sum_{s=3}^m \frac{(2k)^{r+s}}{4rs} \\ &= O\left(\frac{(2k)^{2m}}{m^2}\right). \end{aligned}$$

(c) Finally, by Lemma 3.1, the contribution to (3) from the terms not included in (a) or (b) is bounded above by

$$A \sum_{r=3}^m \sum_{s=r}^m \sum_{l=1}^{r-1} n^{r+s-l-1} \left(\frac{2k}{n}\right)^{r+s-l} \left(1 + \sqrt{\frac{2(l-1)}{n}}\right)^{r+s-2l-2}$$

This expression can be shown to be $O((2k)^{2m}/m^2)$, but we will leave the gory details to the reader. \square

3.3 Theorem Let $m(n) = O(n^{\frac{1}{2}-\epsilon})$ for some $\epsilon > 0$, and suppose that T_1, T_2, \dots is a sequence of positive numbers such that

$$\sum_{i=1}^{\infty} \frac{(2k)^{2m(n_i)}}{m(n_i)^2 T_i^2} < \infty.$$

Then, with probability one, $C(m(n_i), X_i) = O(T_i)$ as $i \rightarrow \infty$.

Proof: This is a simple application of the Borel-Cantelli Lemma and Chebychev's Inequality (see [1] for both) to Theorem 3.2. \square

As a simple application we have:

3.4 Corollary Let $\alpha > 0$ be constant. Then, with probability one, $C(\alpha \log_{2k} n_i, X_i) = O(n_i^{\frac{1}{2}+\alpha})$ as $i \rightarrow \infty$. \square

3.5 Corollary X satisfies Condition (B) with probability one. \square

§4 NOTES

The technique used to prove of Lemma 2.1 can be strengthened considerably, and also used to prove a corresponding upper bound. This will be demonstrated at length in [3], but without the connectivity restriction. For random regular graphs, not necessarily connected, the bound of Theorem 3.2 can be lowered to $O((k-1)^{2m}/m^2)$, which is best possible. See Wormald [4] for many related results.

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