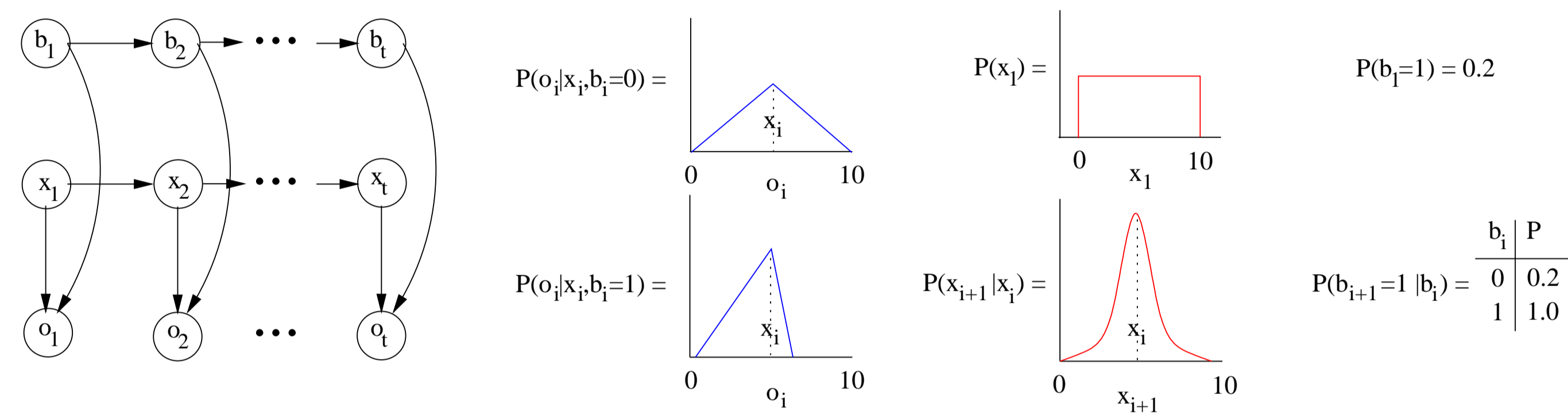


# Symbolic Variable Elimination for Discrete and Continuous Graphical Models

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## Highlight

**Question:** How to do **closed-form exact inference** in discrete and continuous variable (hybrid) **graphical models** with **complex continuous distributions**? E.g.,



**Proposal:** Represent factors by **linear piecewise polynomials**.

**Solution:** Use **symbolic variable elimination** and **extended ADDs (XADDs)**.

## Inference in Hybrid Graphical Models

**Compile:** Bayes nets and MRFs  $\rightarrow$  **Factor Graph** with  $\mathbf{b} \in \{0, 1\}^m$  and  $\mathbf{x} \in \mathbb{R}^n$ :

$$p(\mathbf{b}, \mathbf{x}) = \frac{1}{Z} \prod_{f \in F} \Psi_f(\mathbf{b}_f, \mathbf{x}_f)$$

where each  $\Psi_f(\mathbf{b}_f, \mathbf{x}_f)$  is a **linear piecewise polynomial case statement**, e.g.,

$$\Psi_f(b_1, x_1, x_2) = \begin{cases} b_1 \wedge x_1 > x_2 : & x_1^2 + x_1 x_2 \\ -b_1 \wedge x_1 \leq x_2 \wedge x_1 \geq -x_2 : & x_2^3 \\ -b_1 \wedge x_1 \leq x_2 \wedge x_1 < -x_2 : & x_1^2 x_2^2 \end{cases}$$

**(Conditional) Probability Queries:** for query  $\mathbf{q} = (b_{s+1}, \dots, b_s, x_{t+1}, \dots, x_t)$  and evidence  $\mathbf{e}$  where  $\{\mathbf{b} \cup \mathbf{x}\} \setminus \{\mathbf{q} \cup \mathbf{e}\} = (b_1, \dots, b_s, x_1, \dots, x_t)$  and  $\mathbf{q} = (b_{s+1}, \dots, b_s, x_{t+1}, \dots, x_t)$ :

$$p(\mathbf{q}|\mathbf{e}) = \frac{\sum_{b_1} \dots \sum_{b_s} \int \dots \int_{\mathbb{R}^t} \prod_{f \in F} \Psi_f(\mathbf{b}_f, \mathbf{x}_f) dx_1 \dots dx_t}{\sum_{b_1} \dots \sum_{b_s} \int \dots \int_{\mathbb{R}^t} \prod_{f \in F} \Psi_f(\mathbf{b}_f, \mathbf{x}_f) dx_1 \dots dx_t}$$

**(Conditional) Expectation Queries:** for polynomial function  $Poly(\mathbf{q})$  over  $\mathbf{q}$  (e.g.,  $Poly(\mathbf{q}) = q_1^2 + q_1 q_2 + q_2^2$ ) and evidence  $\mathbf{e}$ , pre-compute  $p(\mathbf{q}|\mathbf{e})$  as above, then:

$$\mathbb{E}[Poly(\mathbf{q})|\mathbf{e}] = \int_{\mathbb{R}^q} [Poly(\mathbf{q}) \cdot p(\mathbf{q}|\mathbf{e})] d\mathbf{q}$$

## Contribution: Symbolic Variable Elimination

**Algorithm 1:** SYMBOLIC VE for conditional probability & expectation queries

**Input:**  $F$ ,  $order$ : a set of factors  $F$ , and a variable  $order$  for elimination

**Output:** a set of factors after eliminating each  $v \in order$

```

begin
  foreach v in order do
    f_v ← 1; F_v ← ∅
    foreach f ∈ F do
      if (f contains v)
        then f_v ← f_v ⊗ f // case multiply: ⊗ (below)
        else F_v ← F_v ∪ {f}
    if (v is boolean)
      then F ← F_v ∪ {f_v=0 ⊕ f_v=1} // case sum: ⊕ (below)
      else F ← F_v ∪ {∫_{v=-∞}^∞ f_v dv} // case definite integral: ∫_{v=-∞}^∞ (below)
  return F
  
```

## Case Operations for Symbolic VE

**Question:** How to do  $\oplus$  (similarly  $\otimes$ ) and  $\max$  (similarly  $\min$ ) on two cases?

$$\begin{cases} \phi_1 : f_1 \\ \phi_2 : f_2 \end{cases} \oplus \begin{cases} \psi_1 : g_1 \\ \psi_2 : g_2 \end{cases} = \begin{cases} \phi_1 \wedge \psi_1 : f_1 + g_1 \\ \phi_1 \wedge \psi_2 : f_1 + g_2 \\ \phi_2 \wedge \psi_1 : f_2 + g_1 \\ \phi_2 \wedge \psi_2 : f_2 + g_2 \end{cases} \quad \max \left( \begin{cases} \phi_1 : f_1 \\ \phi_2 : f_2 \end{cases}, \begin{cases} \psi_1 : g_1 \\ \psi_2 : g_2 \end{cases} \right) = \begin{cases} \phi_1 \wedge \psi_1 \wedge f_1 > g_1 : f_1 \\ \phi_1 \wedge \psi_1 \wedge f_1 \leq g_1 : g_1 \\ \phi_1 \wedge \psi_2 \wedge f_1 > g_2 : f_1 \\ \phi_1 \wedge \psi_2 \wedge f_1 \leq g_2 : g_2 \\ \vdots \\ \vdots \end{cases}$$

**Question:** How to do  $\int_{-\infty}^{\infty} f(x_1) dx_1$  for case statement  $f(x_1)$ ? Expand, swap  $\sum, \int$ :

$$\int_{-\infty}^{\infty} f(x_1) dx_1 = \int_{-\infty}^{\infty} \sum_i \mathbb{I}[\phi_i] \cdot f_i dx_1 = \sum_i \int_{-\infty}^{\infty} \mathbb{I}[\phi_i] \cdot f_i dx_1$$

Hence, we just focus on integrating a single partition:  $\int_{-\infty}^{\infty} \mathbb{I}[\phi_i] \cdot f_i dx_1$ , e.g., let

$$f_i := x_1^2 - x_1 x_2 \\ \phi_i := [x_1 > -1] \wedge [x_1 > x_2 - 1] \wedge [x_1 \leq x_2] \wedge [x_1 \leq x_3 + 1] \wedge [x_2 > 0]$$

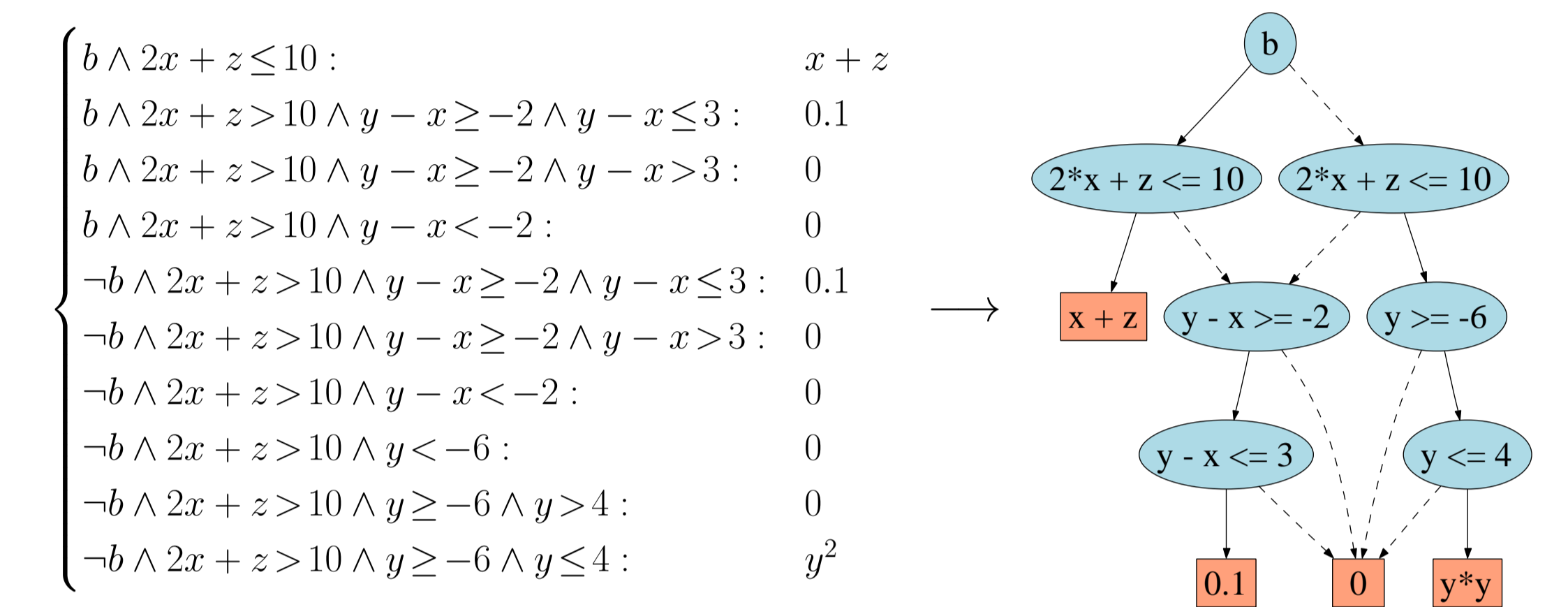
Derive lower/upper bounds  $LB/UB$  from  $\phi_i$  using  $\max/\min$ , integrate, evaluate:

$$LB := \max(x_2 - 1, -1) = \begin{cases} x_2 - 1 > -1 : x_2 - 1 \\ x_2 - 1 \leq -1 : -1 \end{cases} \quad UB := \min(x_2, x_3 + 1) = \begin{cases} x_2 < x_3 + 1 : x_2 \\ x_2 \geq x_3 + 1 : x_3 + 1 \end{cases} \\ \int_{-\infty}^{\infty} \mathbb{I}[\phi_i] \cdot f_i dx_1 = \mathbb{I}[x_2 > 0] \int_{x_1=LB}^{UB} (x_1^2 - x_1 x_2) dx_1 = \mathbb{I}[x_2 > 0] \left[ \frac{1}{3} x_1^3 - \frac{1}{2} x_1^2 x_2 \right]_{x_1=LB}^{UB} \\ = \mathbb{I}[x_2 > 0] \otimes \left[ \left( \frac{1}{3} UB \otimes UB \otimes UB \otimes \frac{1}{2} UB \otimes UB \otimes (x_2) \right) \ominus \left( \frac{1}{3} LB \otimes LB \otimes LB \otimes \frac{1}{2} LB \otimes LB \otimes (x_2) \right) \right] \\ \otimes \mathbb{I}[x_2 > x_2 - 1] \otimes \mathbb{I}[x_2 > -1] \otimes \mathbb{I}[x_3 + 1 > x_2 - 1] \otimes \mathbb{I}[x_3 + 1 > -1] // \text{reason for constraints: enforce } UB > LB$$

## Extended ADD (XADD)

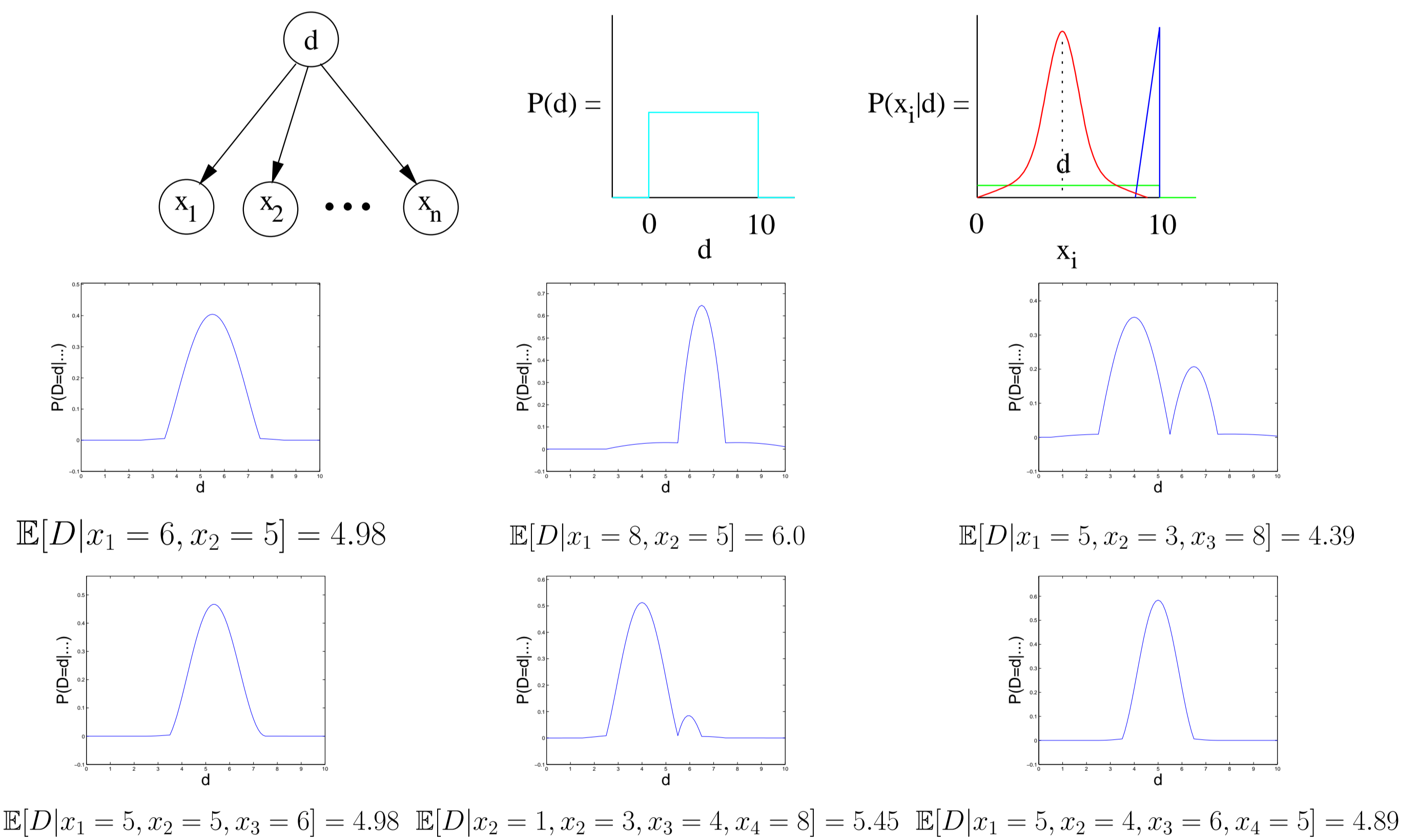
**Question:** How to avoid blow-up in case statements during operations?

**Answer:** Extend algebraic decision diagram (ADD) to represent cases  $\rightarrow$  XADD:

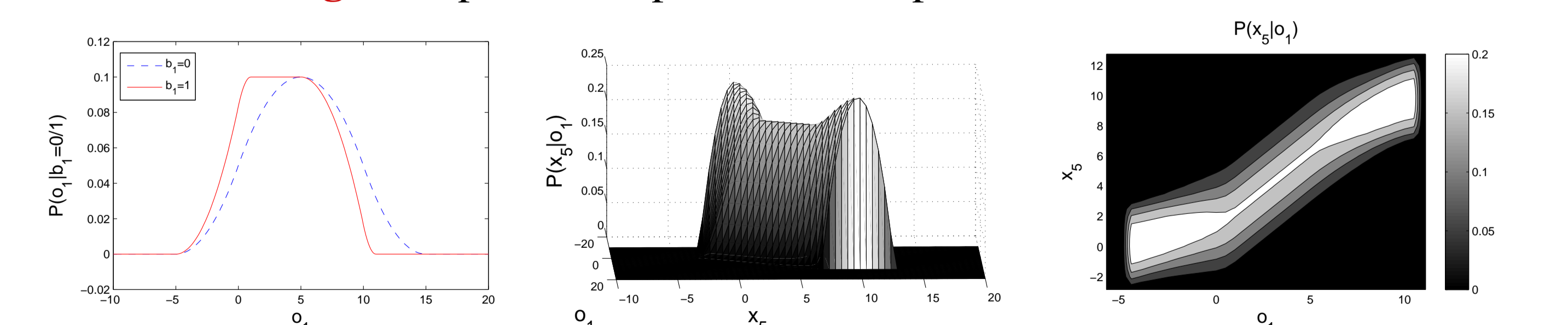


## Example Applications

**Robotics:** Find exact posterior distance given complex sensor noise model:



**Radar Tracking:** Compute exact posterior over position in skewed noise model:



## Conclusion

**First closed-form exact inference for complex continuous graphical models!**