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REALIZATION OF MULTILINEAR AND<sup>1</sup>  
MULTIDECOMPOSABLE MACHINES

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INTRODUCTION

In [1], there appears a category theoretic view of linear systems over a vector space which at the same time extends to a number of other classes of systems, for example, group machines. Our object here is to present some results on multilinear systems in the same vein as [1]. For most of the paper, we shall pose the results in terms of vector spaces. However, at the end of the paper, we indicate generalizations applicable to a category more general than Vect.

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An outline of our results is as follows: we define a K-line multilinear machine in terms of linear, bilinear, ... (K-1)-line multilinear machines via an internal description, and we show that the associated input-output map is multilinear. Then we tackle the converse problem: passage from a prescribed multilinear map to an internal description. We discuss questions of the reachability, observability, and minimality of realizations.

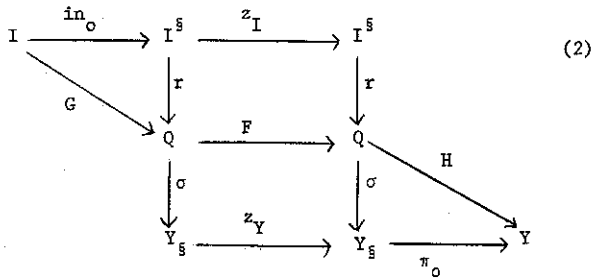
DEFINITION OF A MULTILINEAR SYSTEM

First, we define a linear system as a sextuple  $(Q, I, Y, F, G, H)$  where  $Q, I, Y$  are the state, input and output (vector) spaces, and  $F: Q \rightarrow Q, G: I \rightarrow Q, H: Q \rightarrow Y$  are linear transformations; the system operates according to

$$\begin{aligned} q(k+1) &= Fq(k) + Gu(k) \\ y(k) &= Hq(k) \end{aligned} \tag{1}$$

where  $u, q, y$  are in  $I, Q, Y$  respectively.

In diagram form we have



Here,  $I^S$  denotes the countably infinite copower of  $I$ , which one can identify with left infinite sequences of finite support  $(\dots, 0, \dots, 0, i_k, i_{k-1}, \dots, i_0)$  with the injection  $in_j: I \rightarrow I^S: i_j \mapsto (\dots, 0, \dots, 0, i_j, 0, \dots, 0)$  into the  $(j+1)$ -th slot from the right; one can think of  $i_k$  as  $u(-k)$ . Also,  $Y_S$  denotes the countably infinite power of  $Y$ , which one can identify with right infinite sequences  $(y_0, y_1, y_2, \dots)$  thinking of  $y_k$  as  $y(k+1)$ . One has

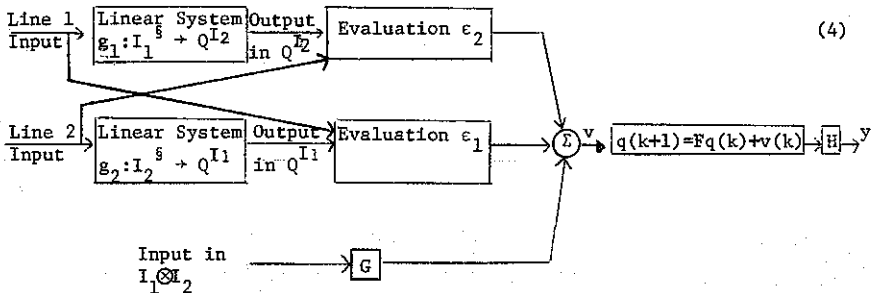
$\pi_j: Y_s \rightarrow Y: (y_0, y_1, y_2, \dots) \rightarrow y_j$ . The morphisms  $z_I$  and  $z_Y$  are shift operators, defined respectively by  $z_I \text{ in }_j = \text{in }_{j+1}$  and  $\pi_{j+1} = \pi_j z_Y$ ;  $z_I$  has the effect of left shift with addition of zero in the zero position and  $z_Y$  of left shift, discarding the element in the zero position prior to the shift. The maps  $r$  and  $\sigma$  are the reachability map (mapping past input sequences into the present state) and the observability map (mapping the present state into a future output sequence, assuming zero future inputs). For a full discussion, see [1].

A multilinear system with  $K$  input lines numbered  $1, 2, \dots, K$  is defined by an equation of the form

$$\begin{aligned}
 q(k+1) &= Fq(k) + \sum_{\epsilon_{i_{j+1} \dots i_K}} \{y_{i_1 \dots i_j}(k) \otimes u_{i_{j+1}}(k) \otimes \dots \otimes u_{i_K}(k)\} + G(u_1(k) \otimes \dots \otimes u_K(k)) \\
 y(k) &= Hq(k)
 \end{aligned}
 \tag{3}$$

where  $I_1, \dots, I_K$  are  $K$  input spaces,  $Q$  is the state space,  $Y$  is the output space;  $u_j \in I_j$ ,  $q \in Q$ ,  $y \in Y$ ;  $y_{i_1 \dots i_j}(k)$  is the output of a known  $j$ -line multilinear system with input spaces  $I_{i_1}, \dots, I_{i_j}$ , the summation is over every partition of  $(1, 2, \dots, K)$  into disjoint nonempty sets  $(i_1, \dots, i_j), (i_{j+1}, \dots, i_K)$ ;  $y_{i_1 \dots i_j}(k) \in \text{Hom}(I_{i_{j+1}} \otimes \dots \otimes I_{i_K}, Q)$  and  $\epsilon_{i_{j+1} \dots i_K}$  is the canonical evaluation map; and  $F, G, H$  are linear maps with obvious domain and codomain.

Figure 1 illustrates the idea for  $K=2$  while Figure 2 illustrates  $K=3$ .



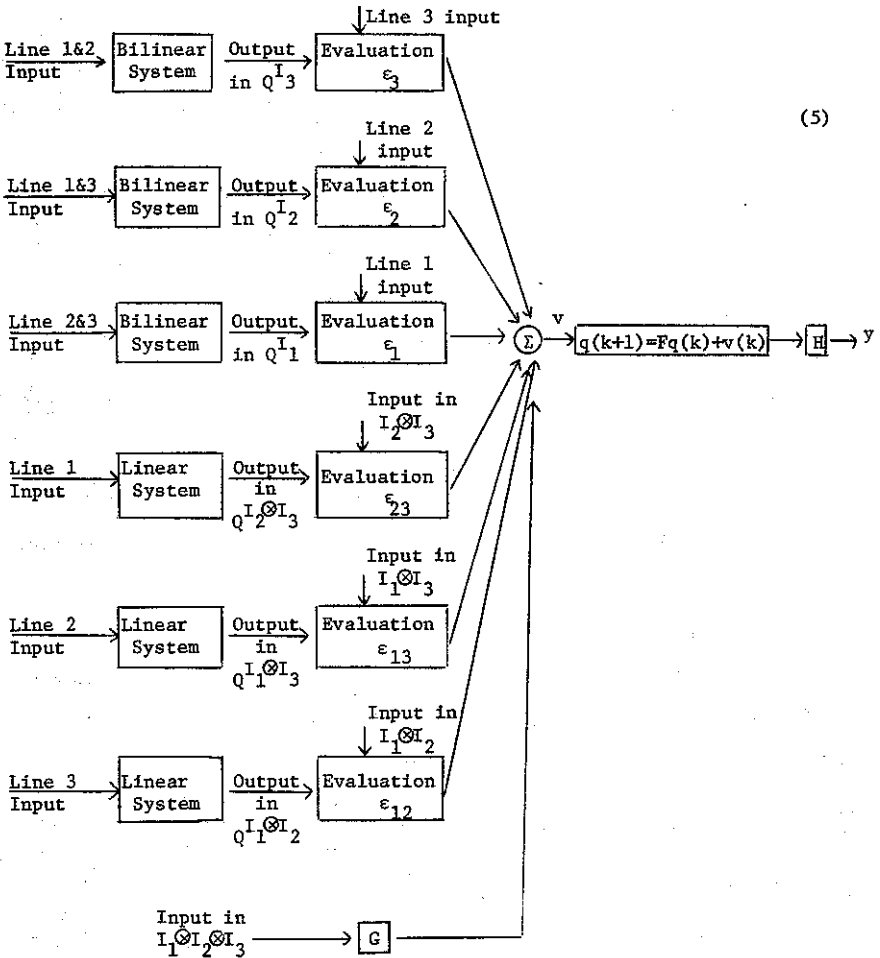


Figure 2

Each bilinear system has its own breakdown into linear systems.

OBTAINING INPUT-OUTPUT DESCRIPTION FROM INTERNAL DESCRIPTION

Theorem 1 A K-line multilinear system defines a linear map  $f: I_1^s \otimes I_2^s \dots \otimes I_K^s \rightarrow Y$  which can be extended to a map  $f^\Delta: I_1^s \otimes \dots \otimes I_K^s \rightarrow Y_s$

such that  $f^\Delta$  possesses the dynamorphic property  $f^\Delta(z_{I_1} \otimes \dots \otimes z_{I_K}) = z_Y f^\Delta$ , and  $\pi_o f^\Delta = f$ .

Remark One can also work with a multilinear map  $\tilde{f} : I_1^s \otimes \dots \otimes I_K^s \rightarrow Y$  which is associated with  $f$  in a canonical way.

In the bilinear case, one defines  $f^\Delta$  as  $\sigma \cdot r$  from the following diagram (showing only one of the linear systems of (4)).

$$\begin{array}{ccccc}
 I_1^s \otimes I_2 & \xrightarrow{I_1^s \otimes \pi_o^2} & I_1^s \otimes I_2^s & \xrightarrow{z_{I_1} \otimes z_{I_2}} & I_1^s \otimes I_2^s \\
 \downarrow g_1 \otimes I_2 & & \downarrow r & & \downarrow r \\
 Q \otimes I_2 & \xrightarrow{\epsilon_2} & Q & \xrightarrow{F} & Q \\
 & & \downarrow \sigma & & \downarrow \sigma \\
 & & Y_s & \xrightarrow{z_Y} & Y_s \\
 & & & & \searrow \pi_o \\
 & & & & Y
 \end{array} \tag{6}$$

The map  $\sigma$  is constructed as for linear systems in [1]. Notice that given the definition (3) is completely equivalent. A set-up like that of (4) has been analyzed in [2-4], with the structure obtained via an analysis of the Nerode equivalence relation for a bilinear map  $\tilde{f} : I_1^s \times I_2^s \rightarrow Y$ . Let us now see how this structure will arise in an apparently quite different way, by generalizing some ideas of [1].

OBTAINING INTERNAL DESCRIPTION FROM INPUT-OUTPUT DESCRIPTION

Suppose  $f : I_1^s \otimes \dots \otimes I_K^s \rightarrow Y$ ; with no attention being paid to reachability or observability. An internal description is obtained as follows: take  $Q = Y_s$ ,  $r = f^\Delta$ ,  $H = \pi_o$ ,  $F = z_Y$

$$\begin{array}{ccc}
 I_1^s \otimes \dots \otimes I_K^s & \xrightarrow{z_{I_1} \otimes \dots \otimes z_{I_K}} & I_1^s \otimes \dots \otimes I_K^s \\
 \downarrow f^\Delta & & \downarrow f^\Delta \\
 Y_s & \xrightarrow{z_Y} & Y_s \xrightarrow{\pi_o} Y
 \end{array} \tag{7}$$

By using the appropriate zero injection,  $I_1^s \otimes \dots \otimes I_j^s \otimes I_{j+1} \otimes \dots \otimes I_K \rightarrow I_1^s \otimes \dots \otimes I_K^s$ , obtain a map  $g_{i_1 \dots i_j}^\bullet : I_{i_1}^s \otimes \dots \otimes I_{i_j}^s \otimes I_{i_{j+1}} \otimes \dots \otimes I_{i_K} \rightarrow Y_s$ . Then  $g_{i_1 \dots i_j}^\bullet : I_1^s \otimes \dots \otimes I_{i_j}^s \rightarrow \text{Hom}(I_{i_{j+1}} \otimes \dots \otimes I_{i_K}, Y_s)$  such that  $\varepsilon_{i_{j+1} \dots i_K} (g_{i_1 \dots i_j}^\bullet \otimes I_{i_{j+1}} \otimes \dots \otimes I_{i_K}) = g_{i_1 \dots i_j}^\bullet$  is uniquely definable.

The mapping  $G$  is defined as composition of  $f^\Delta$  with the injection in  $I_1^s \otimes \dots \otimes I_n^K : I_1 \otimes \dots \otimes I_K \rightarrow I_1^s \otimes \dots \otimes I_K^s$ ; these definitions yield the right commutative diagram from which (by induction on  $K$ ) the internal description can be recovered.

REACHABILITY AND OBSERVABILITY

Definitions of reachability and observability are obtainable via induction on the number of lines of a multilinear system.

Define a reachable multilinear system as one where  $r$  is epi and in case  $K \geq 1$ , the linear, bilinear, ... (K-1)-line multilinear systems within it are also reachable. (Notice that if the second condition fails, the first need not; it is only the input-output properties of the linear, ... (K-1)-line systems which can affect  $r$ , see (4) or (6)).

An observable system is one where  $\sigma$  is mono, and the linear, bilinear... (K-1)-line multilinear systems within it are also observable.

MINIMAL REALIZATION

A minimal realization of a map  $f : I_1^s \otimes \dots \otimes I_K^s \rightarrow Y$  is one which is reachable and observable.

Such a realization can be obtained as follows. Factor  $f^\Delta$  as  $\sigma r$  where  $r$  is epi,  $\sigma$  is mono; take  $Q = r(I_1^s \otimes \dots \otimes I_K^s)$  (Effectively,  $Q = \text{Im}(f^\Delta)$ ). Using

the appropriate zero injection, define  $g_{i_1 \dots i_j}^\bullet : I_{i_1}^s \otimes \dots \otimes I_{i_j}^s \otimes I_{i_{j+1}} \otimes \dots \otimes I_{i_K} \rightarrow \text{Hom}(I_{i_{j+1}} \otimes \dots \otimes I_{i_K}, Q)$ . Induction allows construction of a minimal realization for each such  $g_{i_1 \dots i_j}^\bullet$ .

Define the state-space set  $\mathcal{Q}$  of a K-line multilinear system by:  $Q \in \mathcal{Q}$ , and elements of the state-space set of all linear, bilinear, ... (K-1)-linear multilinear system within the K-line multilinear system are in  $\mathcal{Q}$ . Then we have the expected result:

Theorem The elements, appropriately ordered, of the state-space sets of two minimal realizations of the same  $f = I_1^S \otimes \dots \otimes I_K^S \rightarrow Y$  are isomorphic.

#### CATEGORIES OTHER THAN VECT

Let K be a symmetric closed strict monoidal category with countable co-products and products. One can consider diagrams such as (5) and observe they define, via a generalized sort of recursion, a morphism

$$f = I_1^S \otimes \dots \otimes I_K^S \rightarrow Y, \text{ and one can start with such a morphism and show that}$$

there exist diagrams of the form (5) which define realizations of f. (One

uses the adjointness property of the tensor product to pass between  $\otimes_{i_1 \dots i_j}^S$  and  $g_{i_1 \dots i_j}$  in the construction, which otherwise is the same mutatis mutandis, as for

Vect). If  $\mathcal{E}$ - $\eta$  factorizations exist in K, one can define reachable and observable realizations of a prescribed f. By analogy with [1], we might call these systems multidecomposable.

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